# QMI Lesson 13: Optimization

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The last section discussed the optimization of given functions on given closed intervals. This process can be extended to the optimization of models that simulate real life applications. As in any model, the most difficult part of this process involves the construction of an appropriate model. You may use the following steps to help you create models.

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- 4 Optimize the function *f* over its domain using the methods of Section 4.4.



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- 2 We wish to maximize the area A = lw of the enclosure.
- 3 How do we relate the length to the width? We must do so because we need an equation in a **single** variable.

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By cutting squares out of the corners of a rectangular piece of cardboard and folding up the resulting flaps, the cardboard may be turned into a box (without a lid). If the cardboard is 16 by 10 inches, find the dimensions of the box yielding the largest volume.

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- Well, we have length *I*, height *h*, width *w* and volume *V*. We also may want to draw a figure for this problem since it is a bit complicated.
- 2 We want to optimize volume V = lwh.
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$$V(h) = (16-2h)(10-2h)(h) = 160h-52h^2+4h^3.$$



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- 4 We get, then, that  $V'(h) = 4(40 26h + 3h^2) = 4(3h 20)(h 2)$ . Which means the only critical number in our interval is h = 2. We evaluate V(0) = 0, V(2) = 144, and V(5) = 0, and we get that the volume is maximized when h = 2.

The Metro Transit Authority (MTA) operates a subway line for commuters from a certain suburb to downtown. An average of 6000 people take the train per day, paying \$3.00 per ride. The MTA is considering raising the fare to \$3.50 per ride, but knows that for every \$0.50 increase in fare, an average of 1000 people will choose not to take the train. Show that this increase will reduce revenue, and find the maximum fare increase that will not result in a loss of revenue.

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- 3 How is p related to x? Well, x = 6000 2000(p 3) because 1000 people will leave for every incease of \$0.50 over \$3.00. Thus,  $R(p) = p(12000 2000p) = 12000p 2000p^2 = 2000(6p p^2)$ .

The Metro Transit Authority (MTA) operates a subway line for commuters from a certain suburb to downtown. An average of 6000 people take the train per day, paying \$3.00 per ride. The MTA is considering raising the fare to \$3.50 per ride, but knows that for every \$0.50 increase in fare, an average of 1000 people will choose not to take the train. Show that this increase will reduce revenue, and find the maximum fare increase that will not result in a loss of revenue.

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What's the interval for R(p)? Well, we need  $p \ge 0$  and  $x = 12000 - 2000p \ge 0$ 



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What's the interval for R(p)? Well, we need  $p \ge 0$  and  $x = 12000 - 2000p \ge 0 \implies 0 \le p \le 6$ . So our interval is [0,6]. And R'(p) = 12000 - 4000p. So our critical number is p = 3.

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- 4 Thus, testing R(0) = 0, R(3) = 18000,

- What's the interval for R(p)? Well, we need  $p \ge 0$  and  $x = 12000 2000p \ge 0 \implies 0 \le p \le 6$ . So our interval is [0,6]. And R'(p) = 12000 4000p. So our critical number is p = 3.
- 4 Thus, testing R(0) = 0, R(3) = 18000, and R(6) = 0 gives that the maximum revenue is achieved at p = 3.

- 3 What's the interval for R(p)? Well, we need  $p \ge 0$  and  $x = 12000 2000p \ge 0 \implies 0 \le p \le 6$ . So our interval is [0,6]. And R'(p) = 12000 4000p. So our critical number is p = 3.
- 4 Thus, testing R(0) = 0, R(3) = 18000, and R(6) = 0 gives that the maximum revenue is achieved at p = 3. Thus the maximum fare increase that will not result in a loss of revenue is \$0.00.

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Busch's Baked Beans requires that its cans have a capacity of 54 cubic inches, are cylindrical, and be made of aluminum. Determine the height and radius of the container that requires the least amount of metal, i.e. that has the least surface area.

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Busch's Baked Beans requires that its cans have a capacity of 54 cubic inches, are cylindrical, and be made of aluminum. Determine the height and radius of the container that requires the least amount of metal, i.e. that has the least surface area.

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Busch's Baked Beans requires that its cans have a capacity of 54 cubic inches, are cylindrical, and be made of aluminum. Determine the height and radius of the container that requires the least amount of metal, i.e. that has the least surface area.

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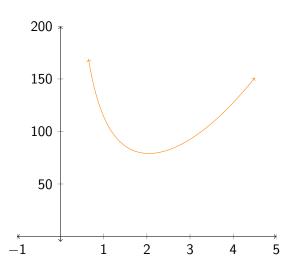
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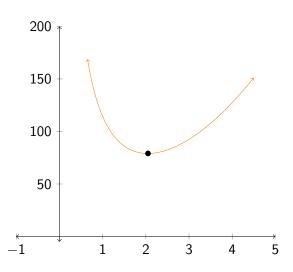
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Notice, S''(r) > 0 for all r in  $[0, \infty)$ , which means that S is concave up on the entire domain considered! Thus,  $r = \frac{3}{\sqrt[3]{\pi}}$  must be an absolute minimum.

# Graph



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A motorcycle retailer sells Shadow 250cc motorcycles exclusively. They estimate that the demand for these motorcycles is 10,000 per year and that they are sold at a uniform rate throughout the year. The cost of ordering a shipment is \$10000 and the yearly cost of storing one motorcycle is \$200. How large should a shipment be and how often should orders be made to minimize ordering plus storage costs?

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$$C(x) = 100x + \frac{10^8}{x}.$$

### Example<sup>1</sup>

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# Assignment

Read 5.1-5.2. Do problems 4, 8, 16, 26, 28, 30, 32, 34 in 4.5.